

# Linear programming



## What is linear programming?

**Linear programming** (LP) finds the **maximum or minimum** of an objective function subject to a set of linear constraints. It is used whenever resources are limited and you want the best possible outcome.

### Key vocabulary:

Term	Meaning
Objective function	The quantity to maximize or minimize (e.g. profit, cost)
Constraints	Linear inequalities restricting the variables
Feasible region	All points satisfying <i>every</i> constraint simultaneously
Corner point	A vertex of the feasible region

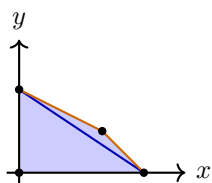
## The fundamental theorem of linear programming

**If an optimal solution exists, it occurs at a *corner point* of the feasible region.**

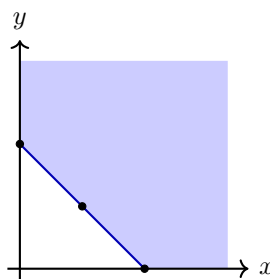
To find the optimum: evaluate the objective function at *every* corner point and pick the best.

**Note:** If the feasible region is **unbounded**, a maximum may not exist. A minimum still exists if all objective function coefficients are positive.

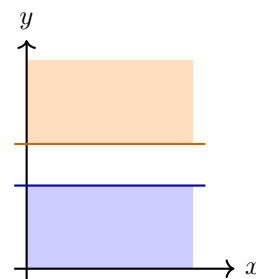
## Solution region types



**Bounded**  
Finite feasible region  
Max & min both exist



**Unbounded**  
Region extends to infinity  
Min may exist; max may not



**No solution**  
Constraints don't overlap  
No feasible region exists

## Step-by-step method

1. **Define variables.** Let  $x = \dots$  and  $y = \dots$  (name the quantities you control).

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2. **Write the objective function.** State what to maximize or minimize:

$$\text{Maximize (or minimize): } P = ax + by$$

3. **Write the constraints.** Translate each restriction into a linear inequality. Include implicit constraints such as  $x \geq 0$ ,  $y \geq 0$  (quantities cannot be negative).

4. **Graph the feasible region.**

- Graph each constraint line (use intercepts).
- Shade each individual half-plane, then identify the *overlap*.
- Use a dashed line for strict inequalities ( $<$ ,  $>$ ); solid line for  $\leq$ ,  $\geq$ .

5. **Find all corner points.** Read coordinates from the graph, or solve the system of two constraint equations at each intersection algebraically.

6. **Evaluate and choose.** Substitute each corner point into the objective function. The largest (maximizing) or smallest (minimizing) value is the answer.

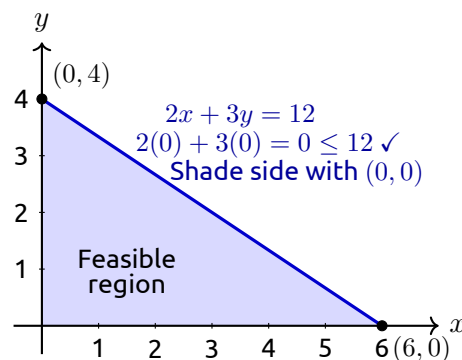
**Tip:** Always check *every* corner point—you cannot tell which is optimal just by looking at the graph.

## Graphing a constraint: quick method

To graph  $ax + by \leq c$  (or  $\geq c$ ):

1. Set  $y = 0$  to get the  $x$ -intercept:  $x = c/a$ .
2. Set  $x = 0$  to get the  $y$ -intercept:  $y = c/b$ .
3. Draw the line through the two intercepts.
4. Test the origin  $(0, 0)$ : if it satisfies the inequality, shade the side containing  $(0, 0)$ ; otherwise shade the opposite side.

**Example:** Graph  $2x + 3y \leq 12$ .



## Finding corner points

Corner points occur where two constraint boundaries **intersect**. Solve the two corresponding equations simultaneously.

**Example:** Find the intersection of  $x + y = 6$  and  $x + 2y = 8$ .

**Solution:** Subtract the first equation from the second:

$$(x + 2y) - (x + y) = 8 - 6 \rightarrow y = 2$$

Substitute back:  $x + 2 = 6 \rightarrow x = 4$ . Corner point:  $(4, 2)$ .

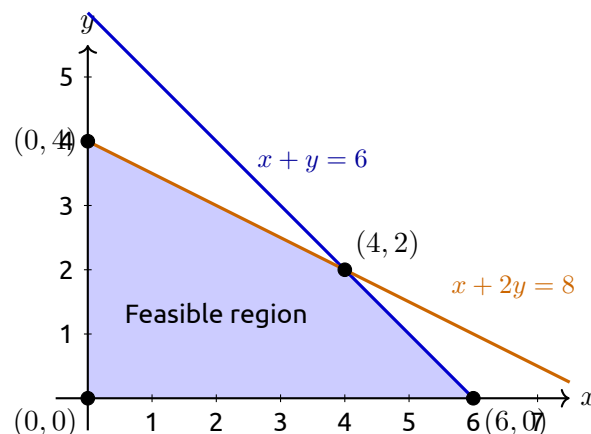
**Note:** Also include the axis intercepts as corner points (where a constraint line meets  $x = 0$  or  $y = 0$ ), and the origin  $(0, 0)$  if it is in the feasible region.

## Full example: maximization

**Example:** Maximize  $P = 60x + 20y$  subject to:

$$x + y \leq 6, \quad x + 2y \leq 8, \quad x \geq 0, \quad y \geq 0$$

**Step 1 – Graph the feasible region.**



**Step 2 – Identify corner points.**

The four corners of the feasible region are  $(0, 0)$ ,  $(6, 0)$ ,  $(4, 2)$ , and  $(0, 4)$ .

The point  $(4, 2)$  is found by solving  $x + y = 6$  and  $x + 2y = 8$  simultaneously (see above).

**Step 3 – Evaluate  $P = 60x + 20y$  at each corner.**

Corner	$x$	$y$	$P = 60x + 20y$
$(0, 0)$	0	0	$60(0) + 20(0) = 0$
$(6, 0)$	6	0	$60(6) + 20(0) = 360$
$(4, 2)$	4	2	$60(4) + 20(2) = 280$
$(0, 4)$	0	4	$60(0) + 20(4) = 80$

**The maximum value is  $P = 360$ , achieved at  $(6, 0)$ .**

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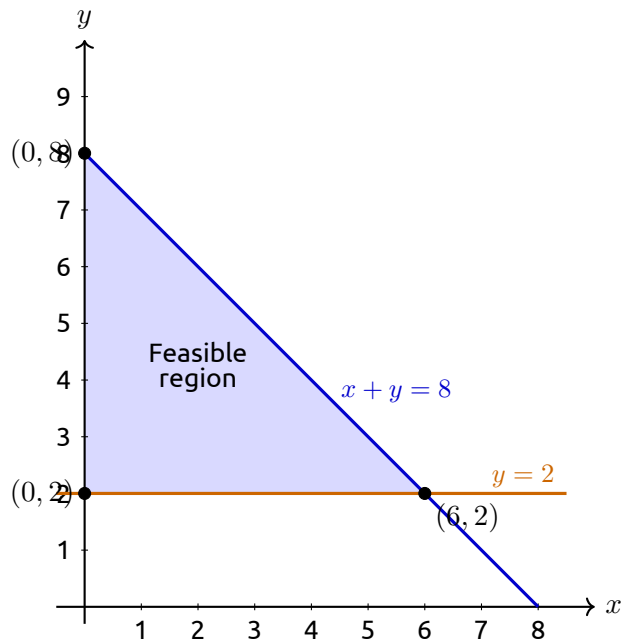
## Full example: minimization

**Example:** A company can produce a **basic** item (profit \$40) and a **premium** item (profit \$60). Assembly capacity limits total units to at most 8, and demand requires at least 2 premium items. How many of each should be made to **minimize cost**, given cost is  $C = 20x + 50y$ ?

Constraints:

$$x + y \leq 8, \quad y \geq 2, \quad x \geq 0, \quad y \geq 0$$

**Corner points:**



Corner points:  $(0, 2)$ ,  $(6, 2)$ ,  $(0, 8)$ .

*Note:*  $(6, 2)$  is found by substituting  $y = 2$  into  $x + y = 8$ , giving  $x = 6$ .

**Evaluate**  $C = 20x + 50y$ :

Corner	$C = 20x + 50y$
$(0, 2)$	$20(0) + 50(2) = 100$
$(6, 2)$	$20(6) + 50(2) = 220$
$(0, 8)$	$20(0) + 50(8) = 400$

**The minimum cost is  $C = \$100$  at  $(0, 2)$ :** produce 0 basic and 2 premium items.

## Common mistakes

**Caution:**

- **Missing a corner point.** Systematically solve every pair of boundary lines to find all intersections – then check each is actually in the feasible region.

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- **Forgetting non-negativity constraints.** Always include  $x \geq 0$  and  $y \geq 0$  unless the problem says otherwise.
- **Testing interior points instead of corners.** The optimal value is always at a corner, never in the interior.
- **Wrong shading.** Test the origin after drawing each line to confirm which side to shade.

## Summary: the LP checklist

- Define  $x$  and  $y$  clearly
- Write objective function: Maximize/Minimize  $P = ax + by$
- List all constraints (include  $x \geq 0, y \geq 0$ )
- Find intercepts for each line; graph and shade correctly
- Identify the feasible region (overlap of all shaded areas)
- Find *all* corner points (axis intercepts + line intersections)
- Evaluate  $P$  at every corner; choose max or min
- State the answer in context

**Practice problems.** For each problem, define variables, write the objective function and constraints, graph the feasible region, find all corner points, and determine the optimal value.

**1.** Maximize  $P = 5x + 4y$  subject to  $x + y \leq 5$ ,  
 $2x + y \leq 8, x \geq 0, y \geq 0$ .

**2.** Maximize  $P = 3x + 5y$  subject to  $x + 2y \leq 10$ ,  
 $3x + y \leq 12, x \geq 0, y \geq 0$ .

**3.** Minimize  $C = 6x + 10y$  subject to  $x + y \geq 6$ ,  
 $2x + y \geq 8, x \geq 0, y \geq 0$ .

**4.** Minimize  $C = 4x + 7y$  subject to  $x + 2y \geq 8$ ,  
 $x + y \geq 5, x \geq 0, y \geq 0$ .

**5.** A company makes chairs (profit \$30) and tables (profit \$70). Each chair takes 2 hrs and each table takes 5 hrs; 40 hrs of labour are available. At most 12 chairs can be made per week. Maximize profit.

**6.** A dietitian needs at least 24 units of vitamin A and 16 units of vitamin B. Food X provides 4 units of A and 2 units of B per serving at \$3. Food Y provides 2 units of A and 4 units of B per serving at \$4. Minimize cost.

**Answers:** 1.  $P = 23$  at  $(3, 2)$    2.  $P = 25$  at  $(2, 4)$    3.  $C = 32$  at  $(2, 4)$    4.  $C = 20$  at  $(5, 0)$

5. Max profit = \$620 at  $(5, 6)$ : 5 tables and 12 chairs   6. Min cost = \$28 at  $(4, 4)$ : 4 servings of each